

Network Traffic Simulation

PASA 96

Real-Time Simulation  
of the  
German Autobahn Network

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*Center for Parallel Computing  
University of Cologne and  
TSASA - Los Alamos National Lab*

## Talk T17

# Real-Time Simulation of the German Autobahn Network

## Topics

- CA–Model of Traffic Simulation
- Network Simulation  
Example: Iterative Routing
- Parallelization
- Outlook and Demo

# Network Traffic Simulation

## Starting Point

# Traffic Simulation

is one of the

**Grand Challenges**

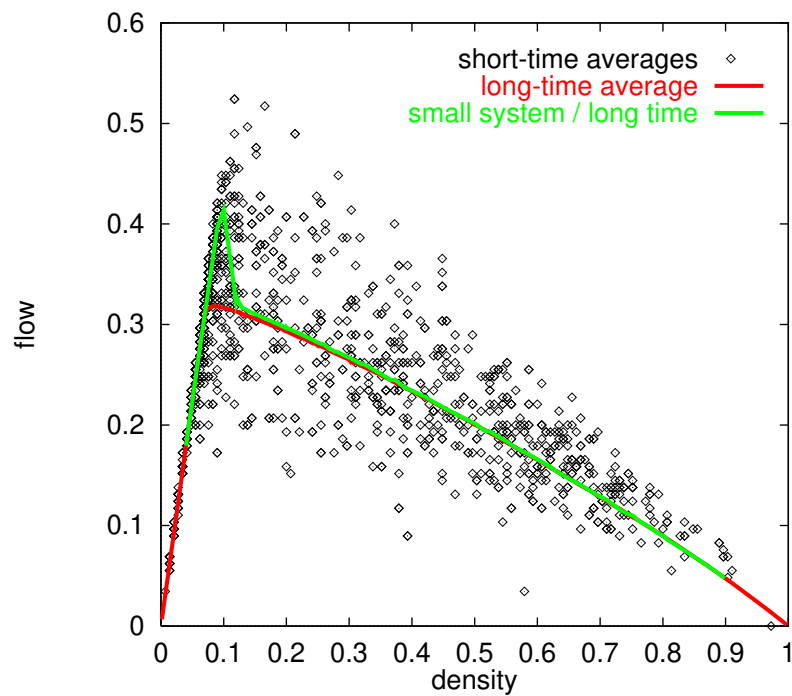
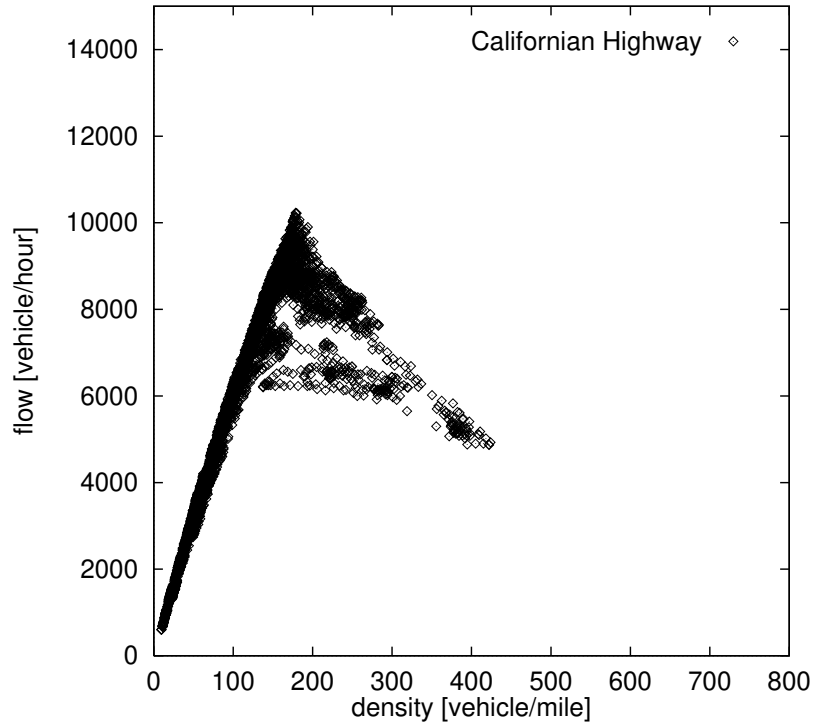
in computer simulation.

## Approach

- Fast algorithms based upon  
→ Cellular Automata (CA)
- Efficient implementation  
→ Parallel Computers

# Network Traffic Simulation

## Motivation



# Network Traffic Simulation

## Work Groups

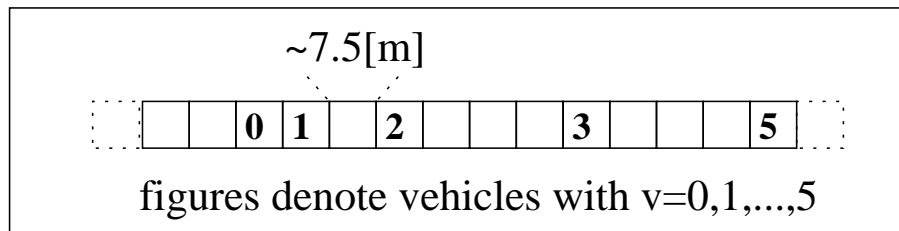
**TRANSIMS** at the **Los Alamos Natl. Lab**  
Los Angeles basin (10 million vehicles)

**ZPR Traffic Group** in **Cologne, Germany**  
*Forschungsverbund Verkehr NRW*  
German Autobahn Network  
(1 million vehicles)

**PARAMICS** in **Edinburgh, Scotland**  
Scottish federal road network

# Network Traffic Simulation

## Single Lane CA Nagel / Schreckenberg (1992)



### 1 Accelerate:

$$v := \min(v_{max}, v + 1)$$

### 2 Avoid crash:

$$v := \min(gap, v)$$

### 3 Randomize:

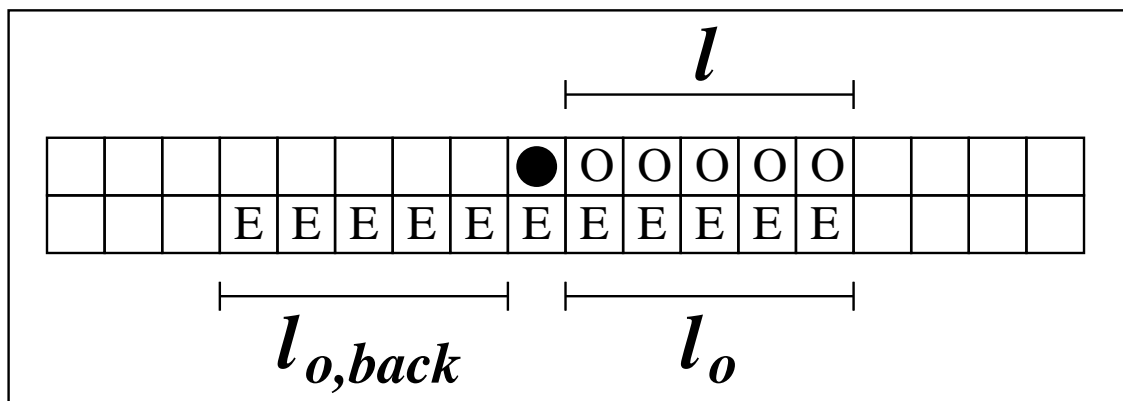
$$rand() < p_{dec} \Rightarrow v := \max(v - 1, 0)$$

Perform Parallel Update

# Network Traffic Simulation

## Lane Changing Rules

- $l$  look ahead same lane
- $l_o$  look ahead other lane
- $l_{o,back}$  look back other lane



### Example

symmetric	asymmetric	
	$L \rightarrow R$	$R \rightarrow L$
$l = v + 1$	no	yes
$l_o = v + 1$	yes	yes
$l_{o,back} = v_{max}$	yes	yes

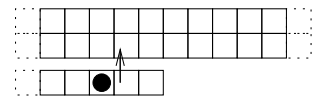
# Network Traffic Simulation

## CA Simulation Network

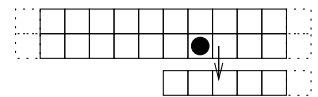
### Building Blocks

- multilane CA

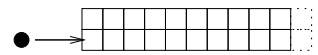
- emission point . . . . .



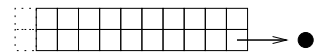
- absorption point . . . . .



- source . . . . .



- sink . . . . .



### Composite Elements

- net terminator (node degree = 1)
- ramp (node degree = 2)
- intersection (node degree = 3,4)

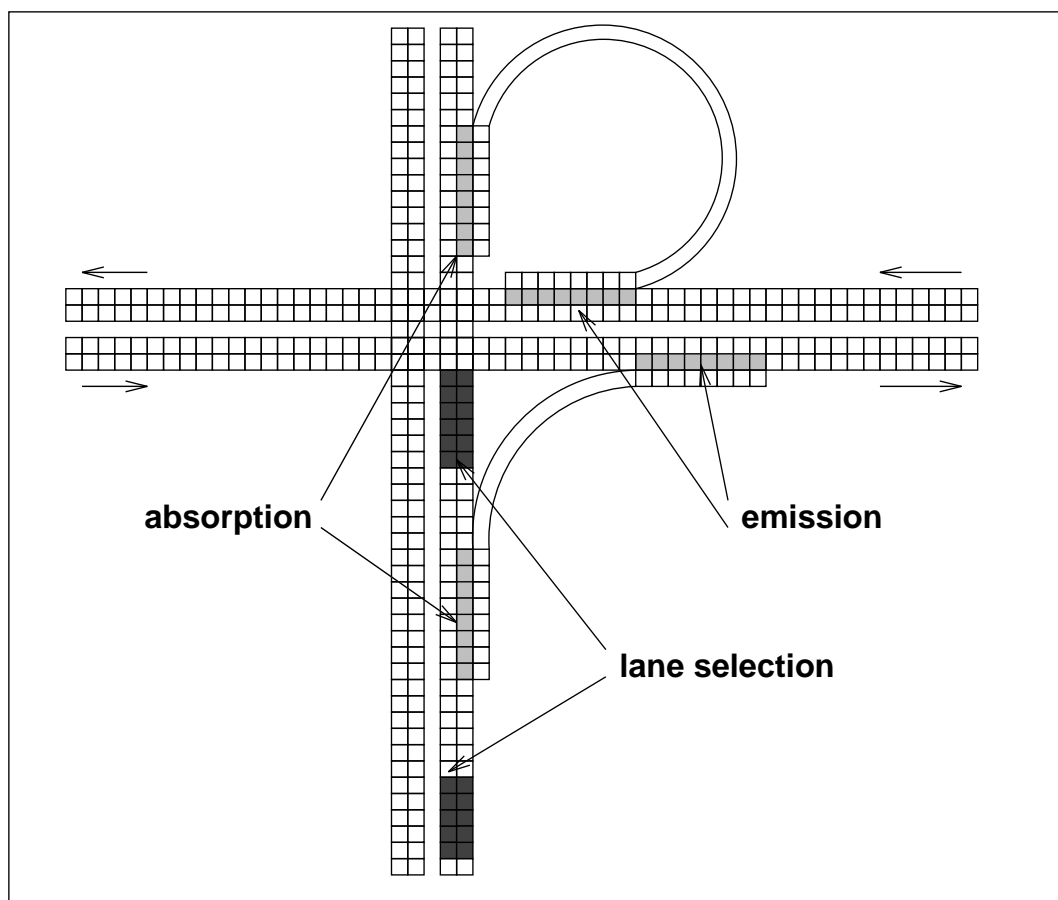


# Network Traffic Simulation

## CA Routeplan Execution

### Vehicles

- behave like 'classical' CA on segments
- have individual route plans
- are absorbed/emitted to follow route



## Example: Iterative Routing

- **Input:** time-dependent origin-destination matrix
- **Output:** consistent set of route plans and edge weights

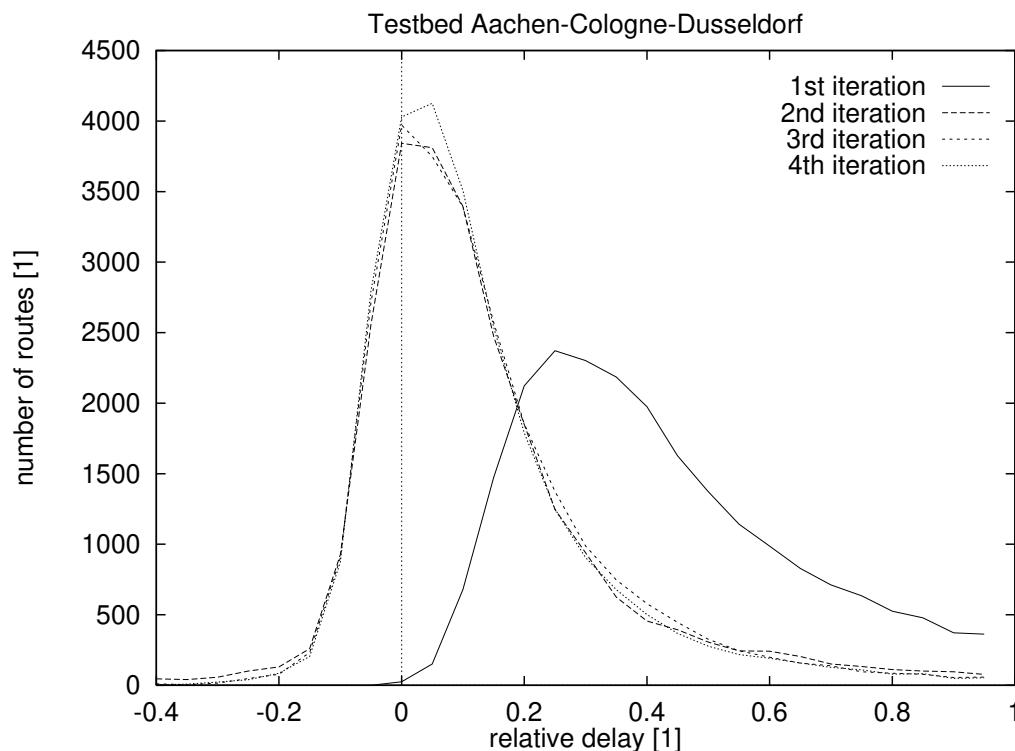
### Iteration of Route Planning

1. preload edge weights  
(e.g. free-flow-velocity)
2. compute route plans for OD-matrix  
(e.g. shortest paths with Dijkstra)
3. simulate route plans while storing actual time-dependent edge weights
4. goto 2

# Network Traffic Simulation

## Routing Example

For low densities ( $\rho = 0.05$ ) the process converges after the first iteration.



### Questions:

Parameter space of convergence?

Quality of prediction?

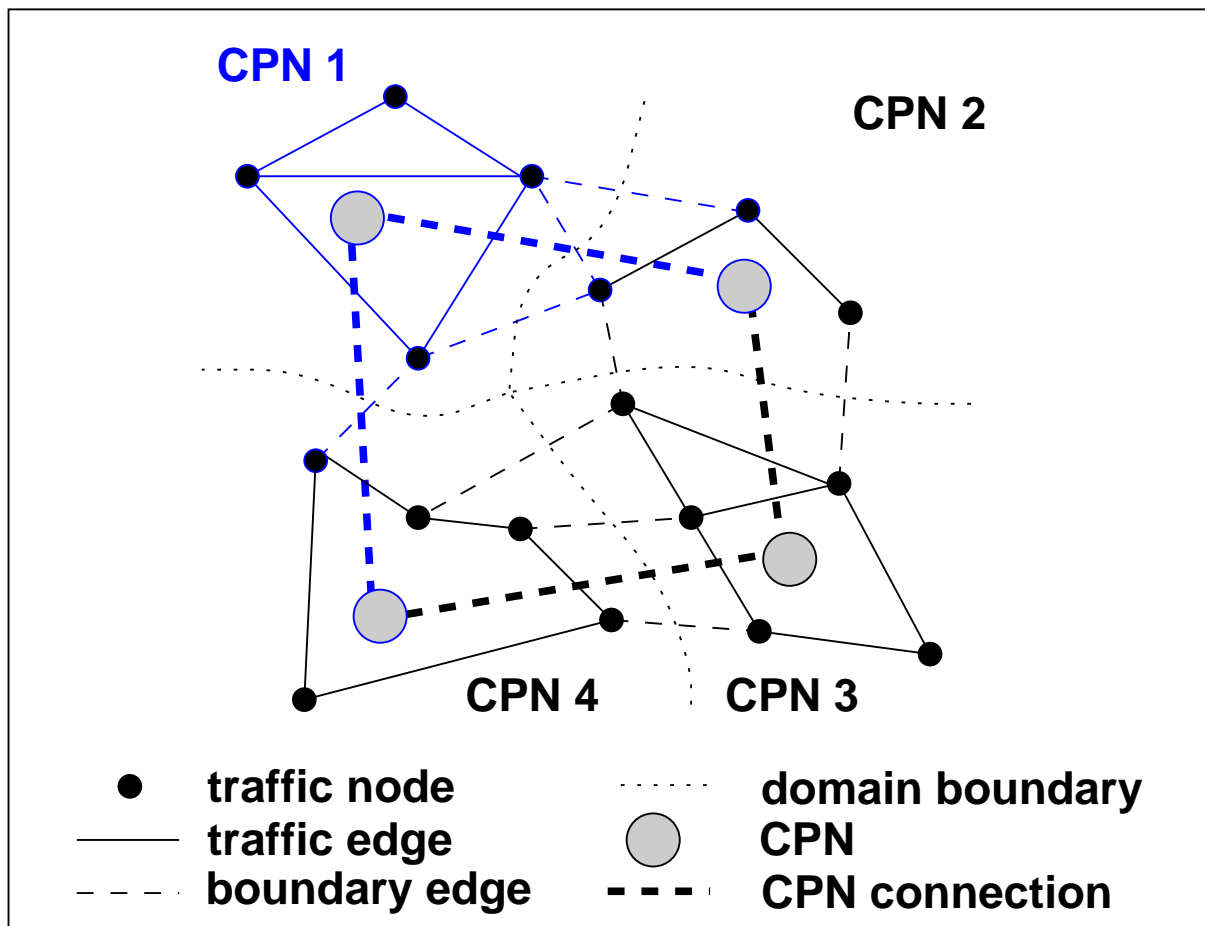
## Parallelization

- traffic network is assigned to a graph of vertices and edges handled by the **Parallel Toolbox**:
  - vertices** correspond to terminators, ramps, and intersections
  - edges** correspond to bidirectional CA multilane segments
- initial **geometric distribution** of vertices (domain decomposition)
- inter-CPN edges handled by **exchange of boundaries**
- dynamic **load balancing**

# Network Traffic Simulation

## Domain Decomposition

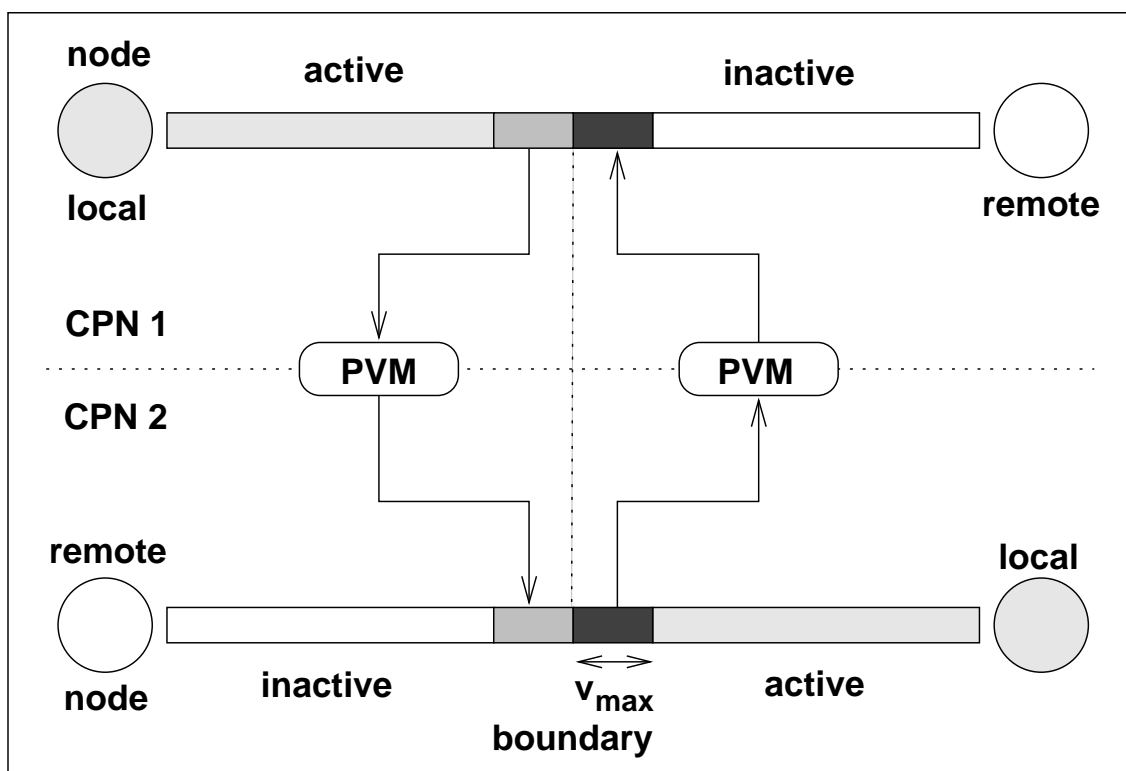
- master CPN has a full copy of inactive network (mainly for graphics)
- each slave CPN has an local active sub network and some inactive dummies



# Network Traffic Simulation

## Boundaries

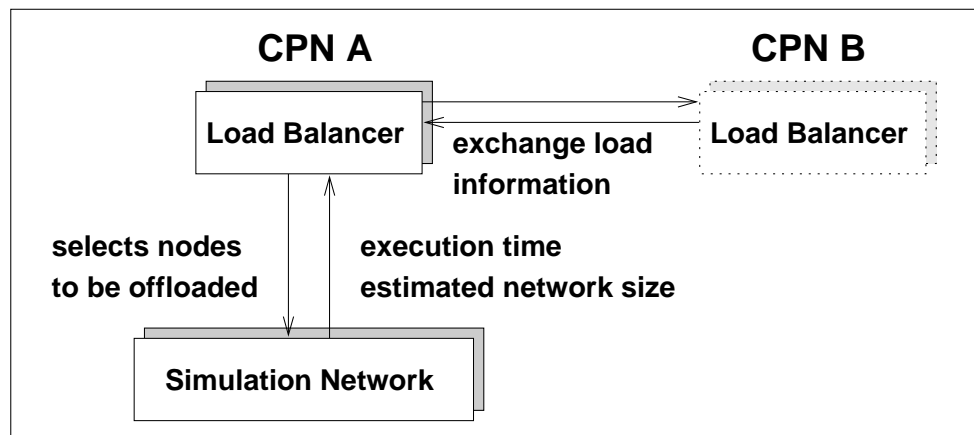
- inter-CPN edges are duplicated with different active ranges
- boundary information is transferred through message passing (PVM)



# Network Traffic Simulation

## Load Balancing(I)

*Local decision, local migration strategy*



Local vertices are offloaded

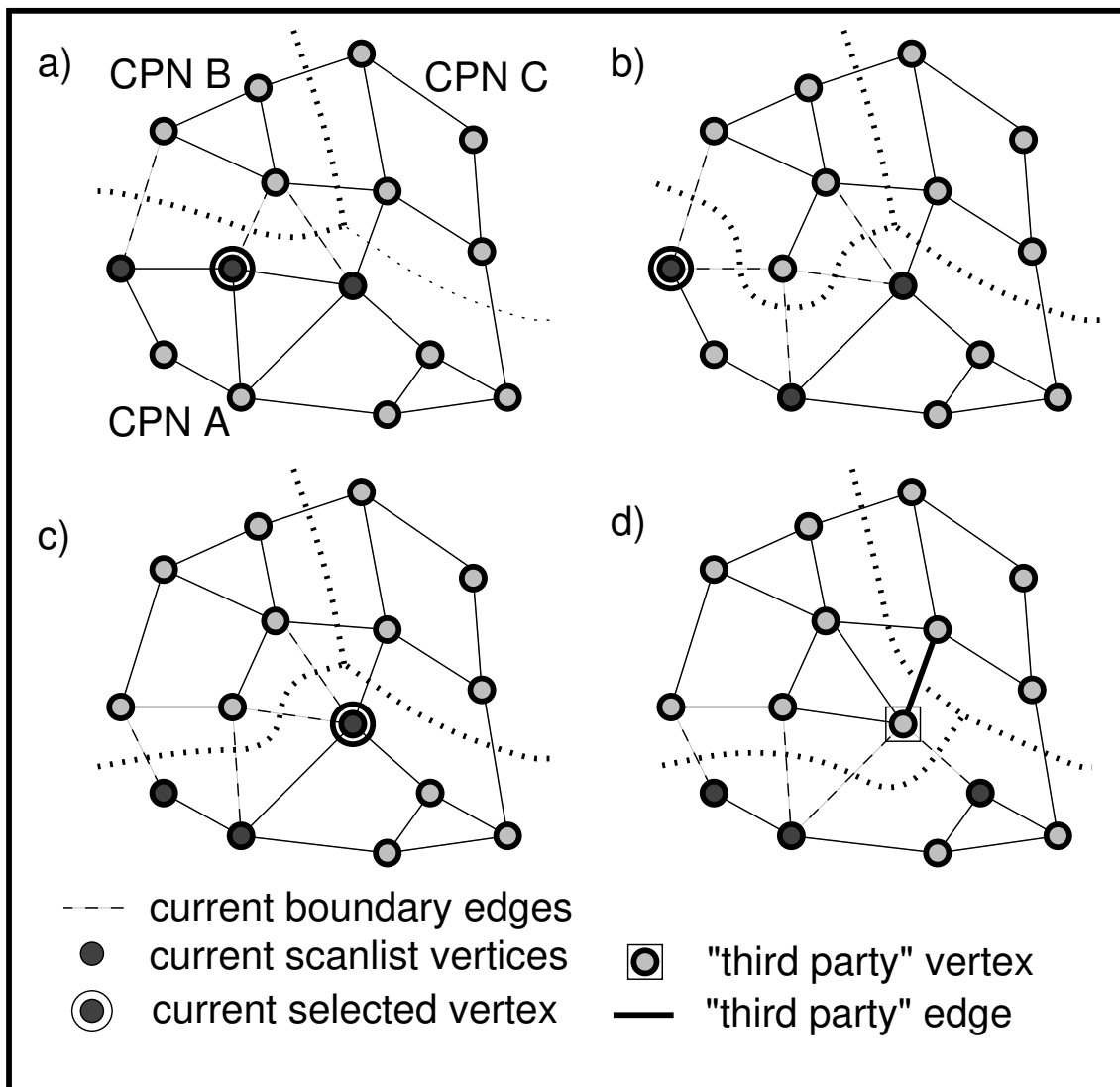
- with **local synchronization** only
- along **common boundaries**
- preferring vertices **furthest from the center** (keep 'nice' shape)
- optionally maintaining **one connected component** per CPN

# Network Traffic Simulation

## Load Balancing (II)

Offloading vertices along the boundaries

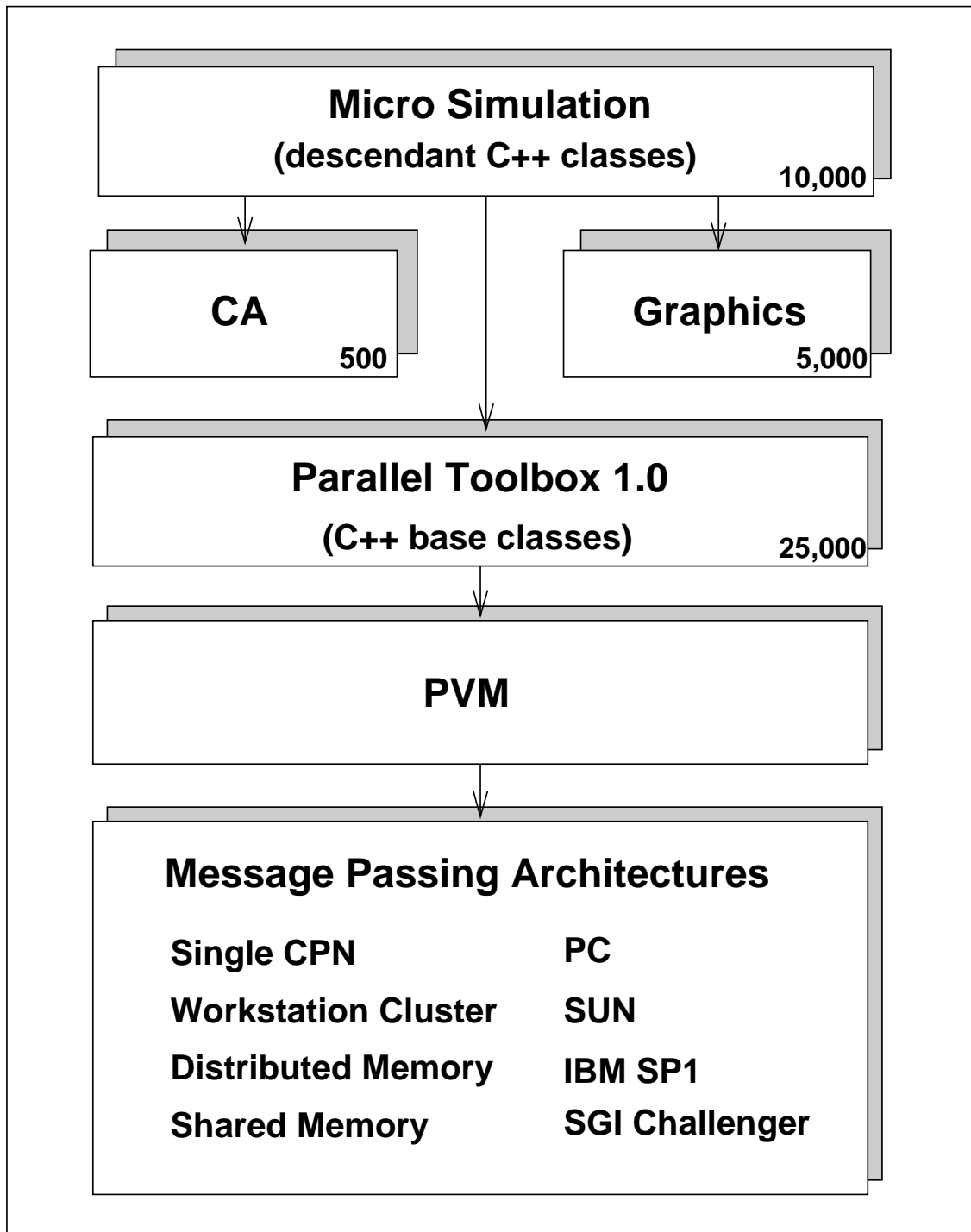
CPN A  $\longrightarrow$  CPN B





# Network Traffic Simulation

## Current Application Structure



# Network Traffic Simulation

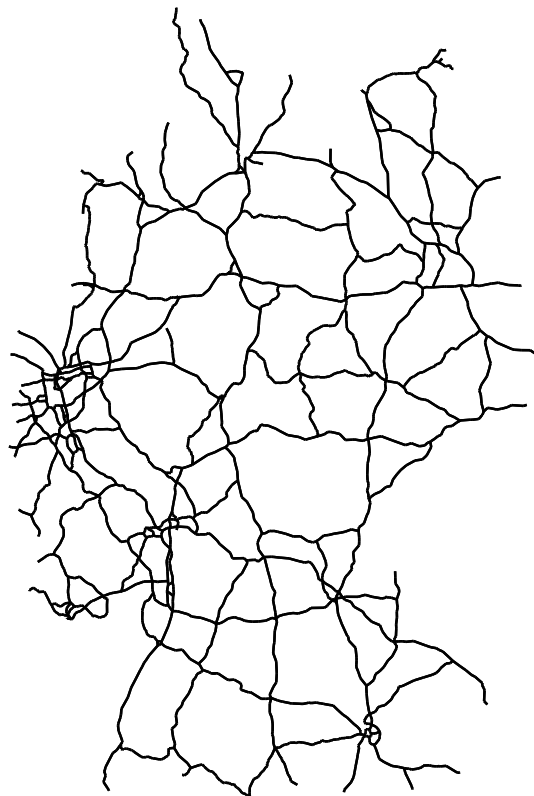
## German Autobahn Network

3300 nodes, 3400 edges

~ 75,000 kilometer (lane corrected)

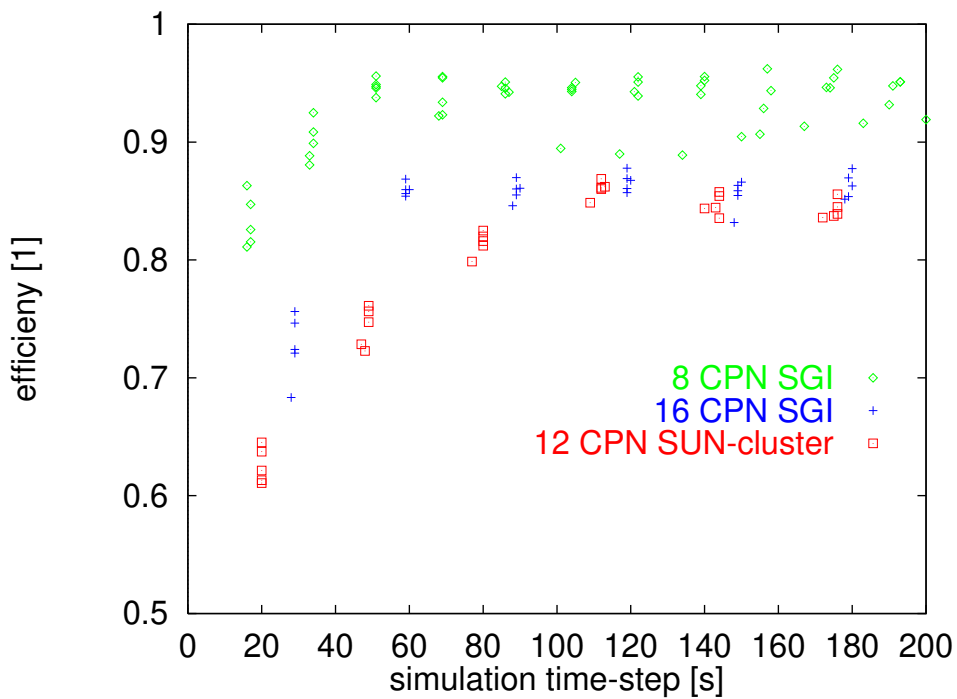
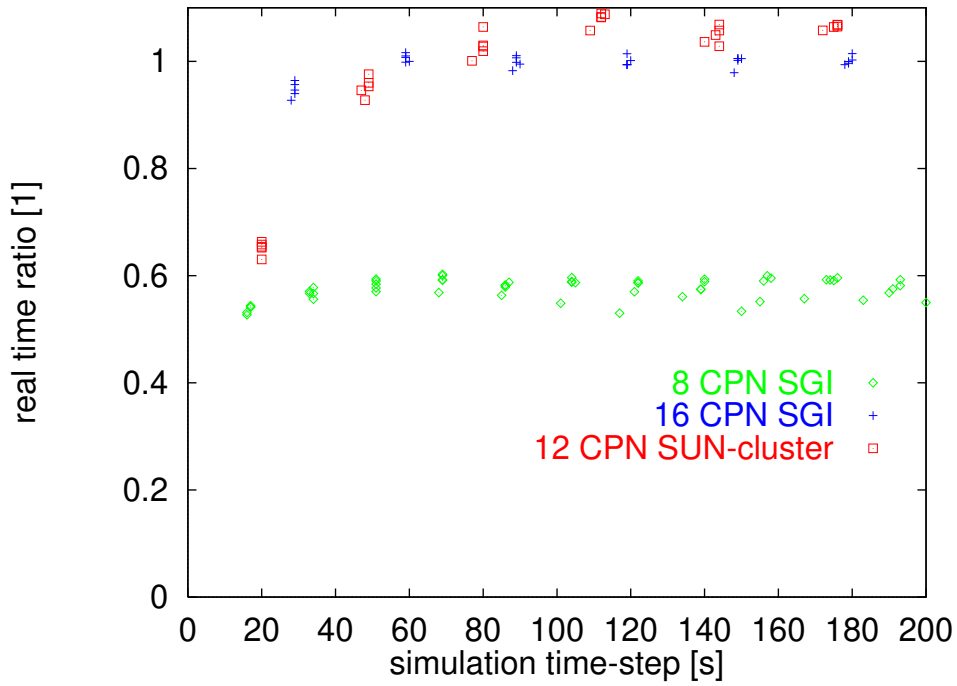
10,000,000 sites

1,000,000 vehicles with routeplans



# Network Traffic Simulation

## Performance for Map FRG



## Efficiency Estimates

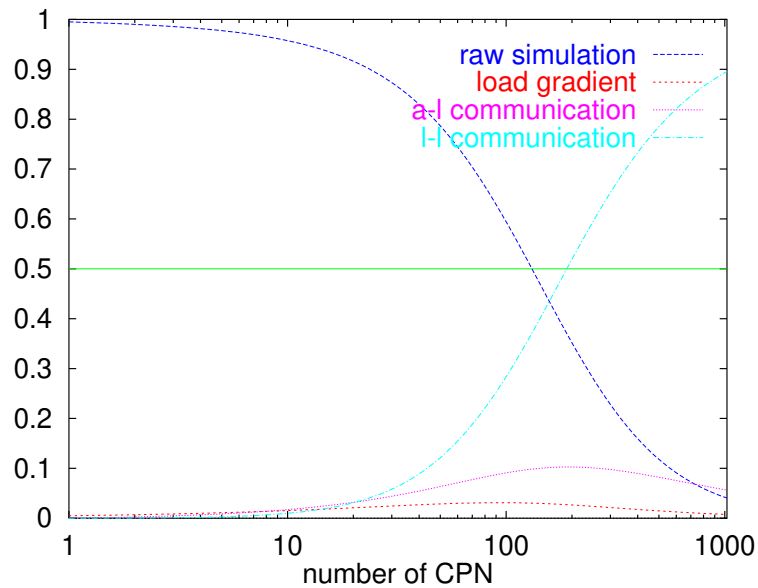
### Criteria

- Application-level communication latency and bandwidth:  
Retrieval, packing, unpacking, and storing of boundary data
- Low-level communication bandwidth:  
Saturation of communication network
- Load gradient:  
Granularity of street network, dynamic load balancing
- Communication topology:  
bus topology or 2-D grid topology

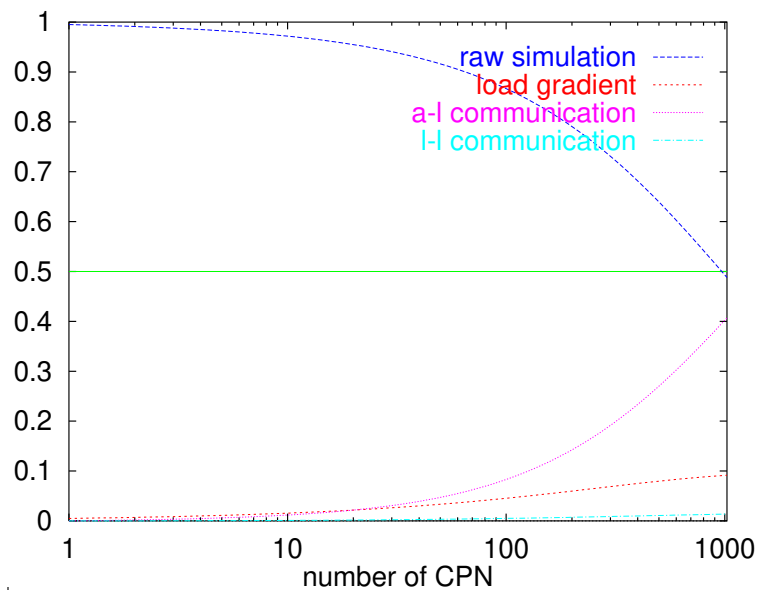
# Network Traffic Simulation

## Extrapolated Efficiency

### SGI Challenger (bus topology)



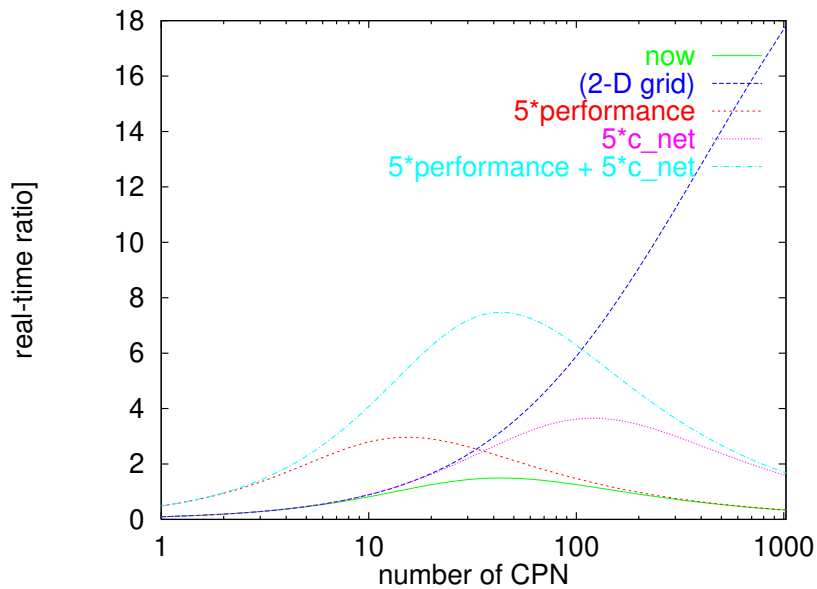
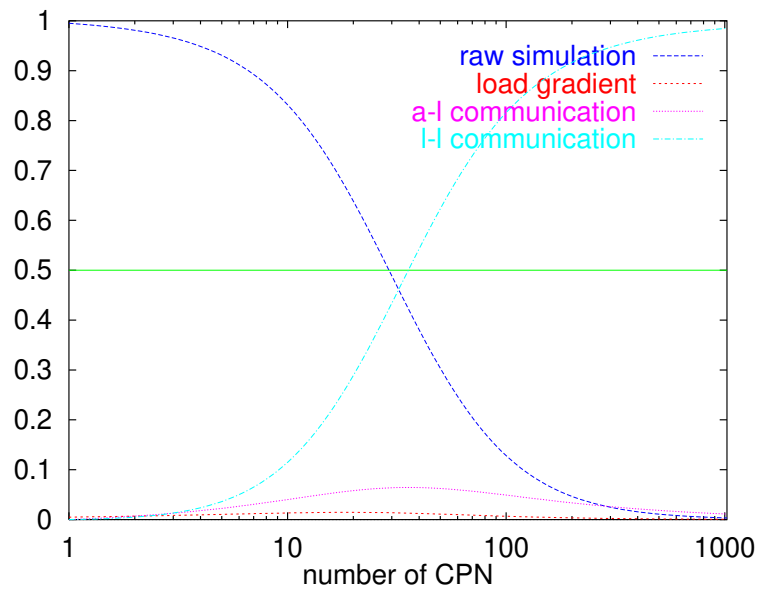
### Intel Paragon (2-D grid topology)



# Network Traffic Simulation

## Achievable Real-time Ratio with Hardware Tuning

SPARC-5 Workstation Cluster (Ethernet)



## Outlook

- Traffic CA

- include vehicle types
- produce more realistic multilane fundamental diagrams
- study net behaviour

- Network Simulation

- online rerouting
- examine stability of routing

- Parallelization

- fewer boundaries
- global corrections of load gradient
- 2-D grid topologies

# Network Traffic Simulation

## Outlook (II)

- Include Shared Memory to allow for hybrid computer networks

